

Applications

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Lines, planes, and hyperplanes

Setup

first welfare theorem

Second welfare theorem

Section 1

Portfolio analysis

Portfolio analysis

- Model of portfolio analysis taken from Simon and Blume
- Good way of reviewing results about matrices and systems of linear equations

Setup

- A assets
- S states of nature
- Initial value of $i = v_i$
- After state revealed, value y_{si}
- Realized return $R_{si} = \frac{y_{si}}{v_i}$
- Matrix of returns $\underbrace{\mathcal{R}}_{S \times A}$

- Choosing budget shares x_i , $\sum_{i=1}^A x_i = 1$ (but individual shares could be positive or negative), call $x = (x_1, \dots, x_A)$ a portfolio

Riskless portfolios

Definition

A portfolio is **riskless** if

$$\sum_{i=1}^A R_{1i} x_i = \dots = \sum_{i=1}^A R_{Si} x_i$$

- Solution to $\mathbf{c} = (c, \dots, c)$

$$\underbrace{\mathcal{R}}_{S \times A} \underbrace{x}_{A \times 1} = \underbrace{\mathbf{c}}_{S \times 1}$$

- Exists if $\mathbf{c} \in \text{Col}(\mathcal{R})$

Insurable states

Definition

A state s^* is **insurable** if \exists portfolio x such that $\sum_{a=1}^A R_{s^*a} x_a > 0$ and $\sum_{a=1}^A R_{sa} x_a = 0$ for all $s \neq s^*$.

- s^* insurable iff \exists solution to

$$\mathcal{R}x = e_{s^*}$$

- So if $e_{s^*} \in \text{Col}(\mathcal{R})$
- All states insurable iff $\text{Col}(\mathcal{R}) = \mathbb{R}^S$

Duplicable portfolios

Definition

A portfolio x is **duplicable** if there is another portfolio w such that $\sum_{i=1}^A x_i = \sum_{i=1}^A w_i$ and

$$\mathcal{R}x = \mathcal{R}w.$$

- i.e. multiple solutions to

$$\begin{pmatrix} \mathcal{R} \\ 1 \dots 1 \end{pmatrix} x = \begin{pmatrix} \mathcal{R} \\ 1 \dots 1 \end{pmatrix} w$$

$$\tilde{\mathcal{R}}x = \tilde{\mathcal{R}}w$$

if and only if $\dim \mathcal{N}(\tilde{\mathcal{R}}) > 0$ if and only if $\text{rank} \tilde{\mathcal{R}} < A$

Section 2

First and second welfare theorems

Welfare theorems

1st: Every competitive equilibrium is Pareto efficient (under some conditions)

2nd: Every Pareto efficient allocation is a competitive equilibrium (under some stronger conditions)

- We will
 - Carefully define competitive equilibrium and Pareto efficiency
 - State and prove theorems

Welfare theorems

- Proofs will use:
 - Vector spaces, linear transformations
 - Should mostly understand
 - Convexity, separating hyperplane theorem
 - Continuity — $\forall \epsilon > 0 \exists \delta > 0$ s.t. ...
 - Maybe difficult to follow, but we will see a lot more of this in calculus
- Economic issues to think about:
 - The result of the two theorems, their assumptions, and how they relate to the real world
 - General role of assumptions in proofs, so can conjecture what happens when assumptions are false

Lines, planes, and hyperplanes

Definition

A **hyperplane** in \mathbb{R}^n is an $n - 1$ dimensional affine subspace. Equivalently, a hyperplane is the set of solutions to a single equation with n variables.

Any hyperplane can be written in the form:

$$H_{\xi,c} = \{x : \langle \xi, x \rangle = c\}$$

where $c \in \mathbb{R}$, $\xi \in \mathbb{R}^n$, and $\|\xi\| = 1$.

Separating hyperplanes

Definition

A set $S \subseteq \mathbb{R}^n$ is convex if $\forall x_1, x_2 \in S$ and $\lambda \in (0, 1)$, we have $x_1\lambda + x_2(1 - \lambda) \in S$.

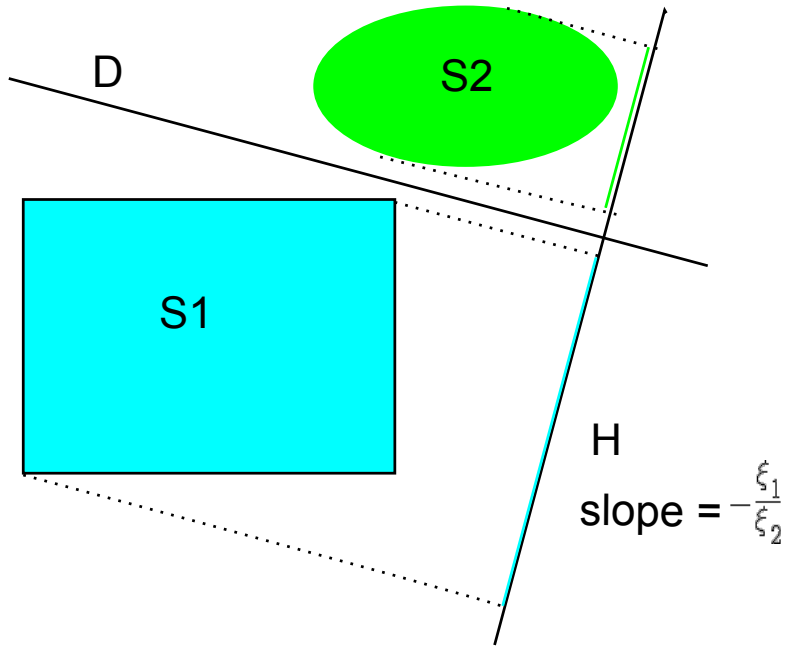
Theorem (Separating hyperplane theorem)

If S_1 and $S_2 \subseteq \mathbb{R}^n$ are convex. If $S_1 \cap S_2 = \emptyset$ then there exists a hyperplane, $H_{\xi c} = \{x : \xi'x = c\}$ such that

$$\langle s_1, \xi \rangle \leq c \leq \langle s_2, \xi \rangle$$

for all $s_1 \in S_1$ and $s_2 \in S_2$. We say that $H_{\xi, c}$ separates S_1 and S_2 .

- Draw picture



Setup

- Commodities $s \in S$, a normed vector space (e.g. S linear subspace of \mathbb{R}^n)
- Consumers $i = 1, \dots, I$ with consumption possibility set $X_i \subseteq S$ (think of as physical constraint, not budget constraint)
- Preference relation \succeq_i s.t.
 - ① (total) $\forall x, z \in X_i$, either $x \succeq_i z$ or $z \succeq_i x$ or both,
 - ② (transitive) $\forall x, w, z \in X_i$, if $x \succeq_i w$ and $w \succeq_i z$ then $x \succeq_i z$,
 - ③ (reflexive) $\forall x \in X_i$, $x \succeq_i x$.

Denote strict preference \succ_i , and indifference \simeq_i

- Could come from utility function i.e. $x \succeq_i z$ iff $u_i(x) \geq u_i(z)$, but does not have to

- Firms $j = 1, \dots, J$ produce y_j from within production possibility sets $Y_j \subseteq S$
 - Example: 3 commodities, food f , clothing c , and labor l . Production function $f = F^f(l)$, and $c = F^c(l)$,

$$Y_j = \{(f, c, l) \in S : l \leq 0 \wedge f \leq F^f(\alpha|l|) \wedge c \leq F^c((1-\alpha)|l|)\} \text{ for } j = 1, \dots, J$$

- $I + J$ tuple $((x_i), (y_j))$ is an **allocation**
- Allocation is feasible if $x_i \in X_i \forall i$, $y_j \in Y_j \forall j$, and the market clears, $\sum_{i=1}^I x_i = \sum_{j=1}^J y_j$.

Definition

An allocation, $((x_i^0), (y_j^0))$, is **Pareto efficient** if it is a feasible and there is no other feasible allocation, $((x_i), (y_j))$, such that $x_i \succeq_i x_i^0$ for all i and $x_i \succ_i x_i^0$ for some i .

Definition

A price system is a continuous linear transformation $p : S \rightarrow \mathbb{R}$. (If S is n -dimensional then p is any $1 \times n$ matrix)

Definition

An allocation, $((x_i^0), (y_j^0))$, along with a price system, p , is a **competitive equilibrium** if

- Ⓒ₁ The allocation is feasible
- Ⓒ₂ (Utility maximization) For each i and $x \in X_i$ if $px \leq px_i^0$ then $x_i^0 \succeq_i x$,
- Ⓒ₃ (Profit maximization) For each j if $y \in Y_j$ then $py \leq py_j^0$

first welfare theorem

Definition

preference relation \succ_j has the **local non-satiation condition** if for each $x \in x_j$ and $\epsilon > 0 \exists x' \in x_j$ such that $\|x - x'\| \leq \epsilon$ and $x' \succ_j x$.

Theorem (first welfare theorem)

if $((x_i^0), (y_i^0))$ and p is a competitive equilibrium and all consumers' preferences have the local non-satiation condition, then $((x_i^0), (y_i^0))$ is pareto efficient.

proof of first welfare theorem

- will show contradiction
- suppose competitive equilibrium not pareto efficient, then \exists a pareto improvement $((x_i), (y_j))$ another feasible allocation s.t.
 - \exists at least one i^* with $x_{i^*} \succ_i x_i^0$
 - the contrapositive of C2 implies $px_{i^*} > px_i^0$
 - all other $x_i \succeq_i x_i^0$
 - local non-satiation and continuity of p imply $px_i \geq px_i^0$
(next slide)

proof of first welfare theorem

- local non-satiation and continuity of p imply $px_i \geq px_i^0$
 - if $px_i < px_i^0$, then local non-satiation implies $\exists x'_i$ very close to x_i such that $x'_i \succ_i x_i \succeq_i x_i^0$.
 - p continuous means that if x'_i is close enough to x_i , then px'_i will be close enough to px_i that $px'_i < px_i^0$ as well.
 - but $x'_i \succ_i x_i$ and $px'_i < px_i^0$ contradicts the definition of competitive equilibrium.
- so in this pareto improvement, we have

$$\sum_{i=1}^i px_i > \sum_{i=1}^i px_i^0 \quad (1)$$

next, we will show this contradicts profit maximization

proof of first welfare theorem

- start with market clearing, apply p , rearrange:

$$\sum_{i=1}^i x_i = \sum_{j=1}^j y_j \quad (2)$$

$$p\left(\sum_{i=1}^i x_i\right) = \sum_{i=1}^i px_i = \sum_{j=1}^j py_j = p\left(\sum_{j=1}^j y_j\right) \quad (3)$$

(4)

- same thing for x_i^0 and y_j^0 with inequality from last slide gives

$$\sum_{j=1}^j py_j > \sum_{j=1}^j py_j^0 \quad (5)$$

but then there is at least on j with $py_j > py_j^0$,
contradicting C3

Second welfare theorem

Definition

A preference relation, \succeq_i , is **convex** if whenever $x \succeq_i z$ and $y \succeq_i z$, then $\lambda x + (1 - \lambda)y \succeq_i z$ for all $\lambda \in [0, 1]$.

Definition

A preference relation, \succeq_i , is **continuous** if for any $x \succ_i z$ there exists a $\delta > 0$ such that for all x' with $\|x - x'\| < \delta$ we have $x' \succ_i z$.

Second welfare theorem

Theorem (Second welfare theorem)

Assume the preferences of each consumer are convex, locally non-satiated, and continuous, and that X_i is convex and non-empty. Also assume that Y_j is convex and non-empty for each firm j .

Suppose $((x_i^e), (y_j^e))$ is a Pareto efficient allocation such that for any price system, p , there is always a cheaper bundle of goods, i.e. $\exists x_i \in X_i$ s.t. $px_i < px_i^e$ for each i . Then there exists a price system, p^e such that $((x_i^e), (y_j^e))$ and p^e is a competitive equilibrium.

Compared to first welfare theorem three more assumptions:

- 1 Convexity of preferences, consumption possibility sets, and production possibility sets
- 2 Continuity of preferences
- 3 Cheaper good condition

Proof of second welfare theorem

- Price system will come from separating hyperplane theorem, first need two disjoint, convex sets
- Let $V_i = \{x \in X_i : x \succ_i x_i^e\}$, set sum of V_i be

$$V = \left\{ \chi \in \mathcal{S} : \chi = \sum_{i=1}^I x_i \text{ for some } x_i \in V_i \right\}$$

V_i convex, so V convex

- Let

$$Y = \left\{ \psi \in \mathcal{S} : \psi = \sum_{i=j}^J y_j \text{ for some } y_j \in Y_j \right\}$$

also convex.

Proof of second welfare theorem

- Show Y and V disjoint
 - Suppose $\chi \in Y \cap V$. Then $\exists x_i \in V_i$ and $y_j \in Y_j$ such that $\chi = \sum_{i=1}^I x_i = \sum_{j=1}^J y_j$, is feasible and $x_i \succ_i x_i^e$
 - Contradicts $((x_i^e), (y_j^e))$ being Pareto efficient
- Separating hyperplane theorem $\Rightarrow \exists p$ s.t.

$$p\chi \geq p\psi \tag{6}$$

- Next, show $((x_i^e), (y_j^e))$ with p is a competitive equilibrium

Proof of second welfare theorem

- Show $p\chi \geq p\chi^e = p\psi^e \geq p\psi$ for all $\chi \in V$, $\psi \in Y$
 - $\psi^e \in Y$, so $p\chi \geq p\psi^e \forall \chi \in V$.
 - $\chi^e = \sum x_i^e = \sum y_j^e = \psi^e$, so

$$p\chi \geq p\chi^e = p\psi^e \quad (7)$$

- Local non-satiation implies we can find $x'_i \succ_i x_i^e$ very close to x_i^e . Continuity of p means that for x'_i close enough we have $|p\chi^e - p\chi'| < \epsilon$ for any $\epsilon > 0$. Therefore $\forall \psi \in Y$

$$p\chi^e + \epsilon > p\chi' \geq p\psi \quad (8)$$

$$\text{so } p\chi^e \geq p\psi$$

Proof of second welfare theorem

- Show $px_i \geq px_i^e$ for each i and all $x_i \in V_i$
 - If not there $px_i + \epsilon < px_i^e$ with $x_i \in V_i$.
 - Local non-satiation and continuity of p means we can choose other $x_k \in V_k$ such that $\left| \sum_{k \neq i} (px_k - px_k^e) \right| < \epsilon/2$
 - But then $p\chi < p\chi^e$ contradicting the previous slide
- Same reasoning but using convexity of Y_j instead of local non-satiation means $py_j \leq y_j^e$ for each j and all $y_j \in Y_j$ i.e. firms are profit maximizing, C3 holds

Proof of second welfare theorem

- Have shown: if $x_i \succ_i x_i^e$ then $px_i \geq px_i^e$
- Need: if $x_i \succ_i x_i^e$ then $px_i > px_i^e$
- Use cheaper good condition
 - If $px_i = px_i^e$ and $\exists x_i' \in X_i$ s.t. $px_i' < px_i^e$, then for all $\lambda \in (0, 1)$,

$$p(\lambda x_i + (1 - \lambda)x_i') < px_i^e \quad (9)$$

- Continuous preferences implies for λ close to 0,

$$\lambda x_i + (1 - \lambda)x_i' \succ_i x_i^e \quad (10)$$

so $\lambda x_i + (1 - \lambda)x_i' \in V_i$, but that contradicts what we have shown.